



Security Enhancement Using NTRU Algorithm Based Cryptosystem for Communication of Classified Information Via Cloud

Aishwarya Rani M R¹, Gururaj Gowda Patil M¹, Mr. Shivanand R D²

¹Student, ²Associate Professor

^{1,2,3}Department of Computer Science and Engineering

Bapuji Institute of Engineering and Technology, VTU, Davanagere, India

ABSTRACT

Distributed storage depends on exceptionally artificial foundation and encourages available boundary point, flexibility and versatility it can oversee profoundly dimensions and tedious information. Henceforth, subcontracting encoded text to a cloud is ended up being a standout amongst the best methodologies for big data stockpiling and admittance. In spite of the fact that gives bounteous security highlights, it is important to verify the client to the most elevated degree with no need of trading off effortlessly utilization and furthermore protected refreshing the encoded text in the cloud in light of another entrance approach as assigned by the information proprietor. These two necessities posture to be a noteworthy test to make the capacity more compelling. There is no real advancement in the territory of admittance approach in the vibrant condition by the conventional frameworks. Now a days, ingress arrangement refresh is critical for improving protection and managing great recurrence of client development. The undertaking goes for executing a protected and obvious admittance organize plot in light of the NTRU encryption algorithm. The deficiency of the current NTRU frameworks will be assessed for corresponding decoding ability and as needs be another NTRU unscrambling calculation will be tried to defeat the unscrambling disappointments of the first NTRU. The plan is exceedingly delicate to the approach refresh it will enables the cloud server to successfully refresh the encoded text when another entrance strategy is indicated by the information proprietor. It likewise empowers (i) the information proprietor and qualified clients to successfully assess the authenticity of a client by check of qualifications and (ii) a client to approve the data gave by different

clients to revise original recuperation. Pre-examination investigation of the plan demonstrates that it will keep qualified clients from swindling and give protection from attacks, for example, the plot assault.

Keywords: *NTRU algorithm, Big data, Cloud, Encryption.*

I. INTRODUCTION

Big Data is a phrase it alludes towards information groups or blends of information collection its dimension, intricacy, and corresponding speed of development lead them hard to caught, overseen, handled or investigated down through traditional knowledge and devices, for example, relational databases and work area insights or representation bundles, inside the occasion important mainly to create them helpful. Difficulties intended for the Big Data incorporate examination, catch, information span, seek, distributing, stockpiling, exchange, perception, questioning, refreshing and data security.

Because of its many sided quality and extensive volume, overseeing Big Data utilizing close by database administration apparatuses is troublesome. A compelling arrangement is mainly to provide subcontract the information toward the cloud server so as to the abilities of putting away Big Data and preparing clients' entrance asks for in a proficient way. For instance, an e-wellbeing applications, the genome data ought to be safely put away in an e-wellbeing cloud because a solitary progression individual genome be approximately 140 gigabytes in an estimate [1]. In any case, while an information

proprietor subcontract associated information in the direction of a cloud, delicate data might be unveiled in light of the fact that the cloud server is not faithful and conviction; along these lines, commonly the encrypted text of the information is put away into the cloud [1]. Be that as it may, how to refresh the encrypted text or cipher text put away into a cloud when another entrance approach is assigned through the statistics proprietor and to confirm the authenticity of a client who expects toward get to the information be at a standstill of extraordinary distress.

A large amount of obtainable methodologies for protecting the outsourced Big Data in mists depend resting on moreover Attribute based encryption (ABE) or mystery sharing. ABE foundational methodologies give the adaptability to an information proprietor to predefinition the arrangement of clients the one is qualified in favor of getting to the information. Mystery distributing instruments enable a secrete on the way to be communal and recreated via firm figure of helpful clients, yet they ordinarily utilize lopsided open key cryptography, for example, RSA for clients' authenticity check, it causes elevated calculation visual projection.

The most testing matter is the means by which to check the authenticity of the clients getting to the subcontracted information in clouds. At present available plans proposed in [1] don't bolster client qualification check. Then again, evident secret distribution construct plans depend with respect to RSA [1] meant for get to authenticity check. Because of various clients require to commonly confirm every other utilizing different RSA activities, equivalent techniques have an elevated statistical visual projection. Moreover, the great asymmetric algorithm for cryptography arrangements, for example, RSA might be out of order through quantum registering sooner rather than later.

The NTRU is an acronym for **Nth degree Truncated polynomial Ring Unit** [9]. The principle trademark be with the intention of amid the encoding and decoding the polynomial duplication is the majority tedious task. This is substantially rapid compare to other deviated cryptosystems, for example, RSA [9]. The NTRU methodology for cryptosystem is a kind of cross section foundational cryptography, and corresponding security depends on top of the briefest vector problem (SVP) in a grid [1]. The significant points of interest of NTRU be portion figuring assault

opposition and illumination of a quick calculation ability.

An enhanced NTRU methodology for cryptosystem (Improved RNS Algorithm) has been projected to defeat the unscrambling disappointments of the first NTRU. At that point a safe and obvious plan in view of the enhanced NTRU and mystery distributing meant for Big Data stockpiling is planned. The cloud server be able to straight forwardly refresh the put away encrypted text or cipher text exclusive of decoding in view of the new access approach determined by the information proprietor, who can approve the restore at the cloud. The anticipated plan be able to confirm the mutual mystery data to keep clients from deceiving and be able to counter different assaults, for example, the intrigue assault. It is additionally esteemed toward be there secure regarding dimension registering assaults because of NTRU.

II. LITERATURE SURVEY

Researchers and specialists posses the chance to modify the range of center by examining the gigantic statistics grouped by the present civilization. To dissect this kind of liberal level informational records, appropriated preparing has been planned as a fiscally keen and reasonable figuring point of view. Regardless, in the view of fact that information passes on confidential data, it ought to be astound starting the cloud and outer aggressors for good, protection, or true blue reasons. Additionally, it has been observed that some expansive level of information examination strategies depend upon second a large amount of key calculation issues, i.e., facilitate variable based math and streamlining and the basic problem and test for flowed handling is the protection of the cloud condition, a broad assortment of rationalities and dimensions have as of late it's been projected by different specialists. Cloud associations suppliers are before long pursuing down the best protection and defense instruments it would create the cloud air protected and ensured intended for their relative clients and it will keep up the complete assurance above the cloud ace affiliation.

Problem Statement

The current Attributed-based encryption (ABE) or mystery distributing frameworks it give verification and adaptability to a information proprietor to predefine the arrangement of clients who are qualified for getting the information however encounter a

misfortune amid its vibrant working that includes standard refreshing of the entrance approach accordingly producing enormous data. The work of asymmetric public key cryptography, for example, RSA intended for client validation causes huge calculation transparency for this situation. The remedy to this problem lies in implementing an algorithm based on a technique that involves regular ingestion policy update, which is highly sensitive to intruder attack and also capable of complex data computation in a dynamic environment. The current Attributed-based encryption (ABE) or mystery distributing frameworks it give verification and adaptability to a information proprietor to predefine the arrangement of clients who are qualified for getting the information however encounter a misfortune amid its vibrant working that includes standard refreshing of the entrance approach accordingly producing enormous data. The work of asymmetric public key cryptography, for example, RSA intended for client validation causes huge calculation transparency for this situation. The remedy to this problem lies in implementing an algorithm based on a technique that involves regular ingestion policy update, which is highly sensitive to intruder attack and also capable of complex data computation in a dynamic environment.

Objectives

The fundamental target of the framework mainly to suggest a protected and irrefutable admittance manage plot intended for the technology Big Data stockpiling located at cloud server and handling the difficulties of the accompanying security administrations:

- To contemplate and break down the present critical admission strategies designed for their deficiencies and pragmatic issues.
- To provide high security, entirely bearing in mind the security of the data storage such that sensitive information is not vulnerable.
- To demonstrate that proposed scheme can resist various attacks such as the collusion attack via a rigorous analysis.

III.METHODOLOGY

Key, individuals who coordinate through dispatcher be able to refresh sender's information on the cloud server. When sender will distribute a group S of its information with collector, it can figure the total key KS for recipient through the stage of fetching function is similar to Extract (MSK, S). Because KS is only a

steady dimension key, it is anything but difficult near be sent to beneficiary by means of a protected email.

IV.ARCHITECTURE

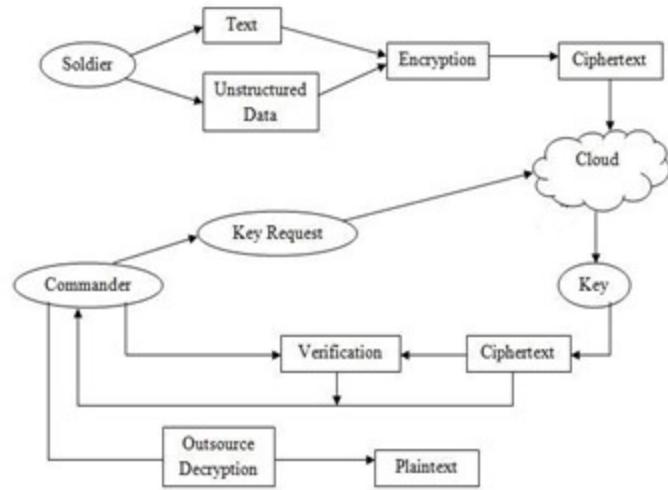


Figure 1: Secure and Verifiable access control methodology

The figure 1 illustrates the secure and verifiable access control methodology. Planned arrangement scheme consists of three polynomial-time algorithms as follows:

- Key Generation Step
- Encryption Step
- Decryption Step
- The information proprietor sets up the general population framework arguments by means of association and creates an open or master security key combine through Key production. Communication knows how to encode by means of Encrypt by any individual who likewise chooses what figure content class is related through the normal text communication to be scrambled.
- The information proprietor be able to utilize the master security key to produce a total decoding key for an arrangement of encrypted cipher text classes through Extract. The created or produced security keys be able to share with entrust safely (by means of protected messages or safe and sound gadgets).
- In conclusion, some client by means of a total key it will be decrypted by any cipher text gave by the cipher text medium is enclosed in the total key by means of Decrypt.
- Assume dispatcher needs to distribute his information m_1, m_2, \dots, m_i to the cloud server. To begin with Setup is performed to get arguments

and carry out Key Generation stage mainly to receive people in public or master security key match (PK means of primary key; MSK means of master security key).

- The scrambled information's are transferred intended to the server. By means of param and PK it says that primary.

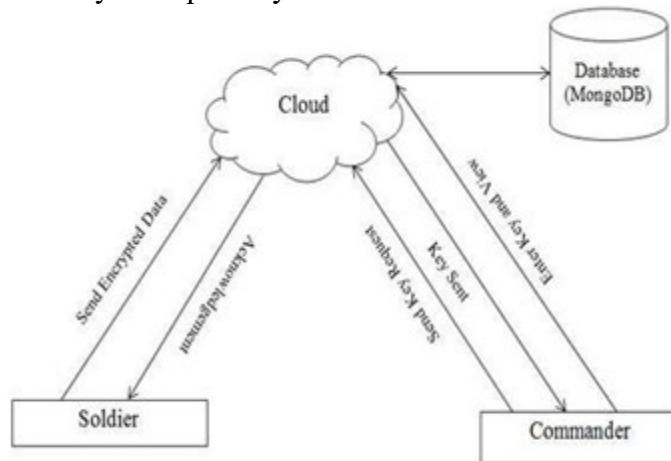


Figure 2: System Architecture

The secure and verifiable access control scheme cryptosystem uses the above system architecture. Application makes use of MongoDB to store and process the data. The frontend of the application is designed using servlet. The proposed system uses NTRU (N^{th} degree Truncated polynomial Ring Unit) algorithm to encrypt the data and decrypt it using key generated by the NTRU algorithm. Thus the application allows only intended users to get verified access to important data.

Step1: Soldier will upload the data into cloud. After uploading data successfully commander will get the intimation mail.

Step2: Commander requests the key from cloud. Key is generated using NTRU algorithm. Once he get the key from cloud he can view the original data.

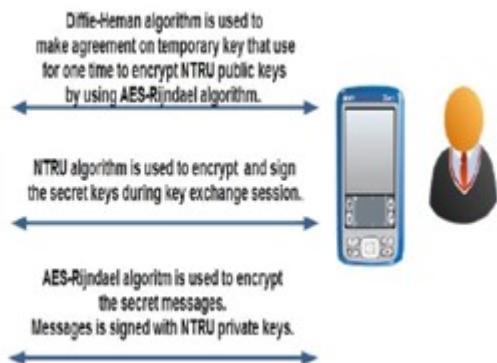
Step3: Commander will reply to soldier.

Step4: Soldier will view the reply sent by commander.

V. ALGORITHM

NTRU is an acronym it illustrates N^{th} degree shortened function ring. The important features are that for the duration of the encoding and decoding the function reproduction is the majority tedious complex process. It is a great quicker than additional asymmetric encryption algorithm, for example RSA asymmetric algorithm and elliptic curve encryption algorithm. The NTRU was developed during 1996 by

scientists Jeffrey Hoffstein, Joseph H. Silverman, and Jill Pipher. Afterward finally during 1996 scientists besides with Daniel Lieman set up the NTRU encryption algorithm. The scientists were measured on increasing up the procedure.



NTRU Keys and Parameters

- N - The polynomials in the ring R have degree $N-1$.
- q - The large modulus to which each coefficient is reduced.
- p - The small modulus to which each coefficient is reduced.
- f - A polynomial that is the private key.
- g - A polynomial that is used to generate the public key h from f (Secret but discarded after initial use)
- h - The public key, also a polynomial
- r - The random “blinding” polynomial (Secret but discarded after initial use)
- d - Coefficient

Key Generation Step

- **Step 1:** User B randomly chooses 2 small polynomials f and g in the R
- **Step 2:** The inverse of f modulo q and the inverse of f modulo p will be computed

$$f * f_q^{-1} = 1 \pmod{q}$$

$$f * f_p^{-1} = 1 \pmod{p}$$
- **Step 3:** Product of polynomials will be computed:

$$h = p * ((Fq)^*g) \pmod{q}$$

Encryption Step

- **Step 1:** User A has a message to transmit.
- **Step 2:** User A puts the message in the form of polynomial m whose coefficients is chosen modulo p between $-p/2$ and $p/2$.
- **Step 3:** Randomly chooses another small polynomial r.
- **Step 4:** Computes the encrypted message:

$$e = r * h + m \pmod{q}$$

Decryption Step

- **Step 1:** Client B obtains a memo e commencing from A and wants to decode it.
- **Step 2:** By using confidential function polynomial f. Client B calculates a function.
- **Step 3:** Client B requires to select constants of a that resides in an intermediate of span q.
- **Step 4:** Client B calculates the function $b = a \pmod p$. Client B decreases each and every constants of a mod p.
- **Step 5:** B utilizes the additional confidential function f_p to calculate $c = f_p * b \pmod p$, it is the unique memo of A.

VI. IMPLEMENTATION DETAILS

Software Requirements

- | | |
|-------------------------|------------------------|
| ➤ Operating System | : Windows 7 and Above |
| ➤ Web Technology | : Servlets, JSP |
| ➤ IDE | : Eclipse Mars |
| ➤ Application server | : Apache Tomcat 8.0 |
| ➤ Hadoop Database | : MongoDB |
| ➤ Database Connectivity | : Robomongo-0.8.5-i386 |

Hardware Requirements

- | | |
|-------------|----------------------|
| ➤ Processor | : Intel I3 and Above |
| ➤ RAM | : 4 GB |
| ➤ Hard Disk | : 500 GB |

A. Procedure to upload data

The below described procedure allows soldier to upload data related to current situation. All data will be successfully stored into „report“ collection. After the data is successfully uploaded corresponding soldier will be acknowledged.

Input: Data consist of soldier name, communicating mail id, confidential secret id, place, date, document and state.

Process: Uploading soldier information. Information will be saved in report collection.

Create collection db object;
Construct a mongo client object;
Fetch the database named "secure data";
From database get collection report;
construct basic database object;

Append soldier information to database such as name, mail id, confidential id, date, state details;

Output: Information uploaded effectively message displayed on the screen.

B. Procedure for data Encryption

The below procedure describes how the uploaded data is encrypted using public key. The encrypted data is stored in „plan“ collection.

Input: Information on the way to be encoded. i.e., database entities.

Process: Encoding the information via utilizing NTRU information encoding step.

Construct a multipart request object y providing parameters name, directory name; Fetch parameter names;

```
for (var param in parameters)
  {fetch next parameter name;
  if (parameter name == "detail")
    {e = Get multi parameter name;
    Fetch the confidential key;
    Get bytes and fetch AES key;
    Construct encoded text nothing but
    cipher text;
    Initialize cipher text;
    Get bytes associated with the
    cipher text;
    Generates the encrypted text;}
```

Output: Information will be encoded effectively and saves in database.

C. Procedure for Key Generation

The below pseudocode describes the key generation procedure and the generated key is sent to commander mail-id through SMTP protocol. The generated key will be stored in key collection.

Input: Press waiting link button;

Process: Produce the Key

- Generate an integer y random number;
- Convert generated random into to string;
- Set attribute value to mail id;
- create boolean variable session debug and
- assign value to true ;
- Fetch system properties;
- Put host properties values;
- Assign transport protocol value to SMTP; Assign authentication value to true; Assign mail smtp value to true;
- Assign port value as 465; Assign fallback value to "false"; Assign class value to

- SSL_FACTORY; Construct a mails session object with props; Assign set Debug to session Debug; Create a new mime object;
- Assign internet address to new mime object;
 - Construct internet address array;
 - Fetch message recipients;
 - Set subject value to msg object; Set content value to msg object; Construct transport object and set value to smtp;
 - Connect transport object by using corresponding user name and password;

Output: Key production will be effectively done and saved in database.

D. Procedure for data Decryption

The below procedure describes the data decryption module. Commander decrypts the data by using generated key to view original data.

Input: Cipher text will be given as an input;

Process: Decoding the information through utilizing NTRU information decoding step

- Fetch a secrete id value from the request parameter;
- Fetch a key value from the request parameter;
- Construct a collection object from the database collection;
- Construct a mongo client object and initialize it with a values;
- Get a mongo client database with a name secure data;
- Fetch user name from the database collection;
- Construct a basic database object with parameters secrete id; Construct a query object;
- Assign secrete id value to query;
- Assign id value to query;
- Search collection to find the query;

While (until cursor has next)

{display the response from viewport jsp page;}

Display the information saying (entered key is wrong);

Output: Information decoded effectively.

VII. RESULTS

The product is demonstrated module-wise with respect to the users such as soldier and commander.

The screen shots given below depict the project results.

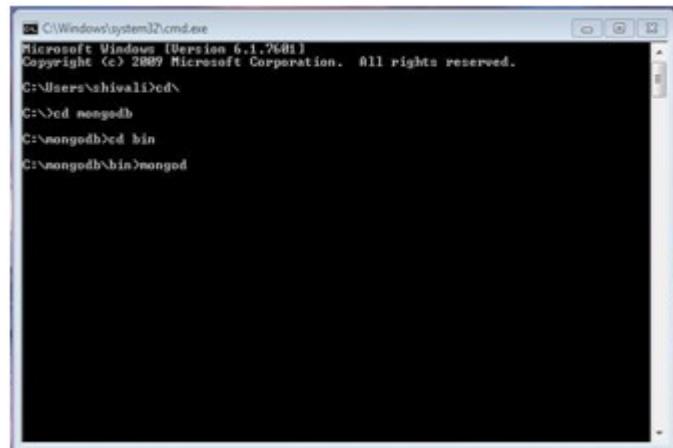


Figure3: Launching MongoDB

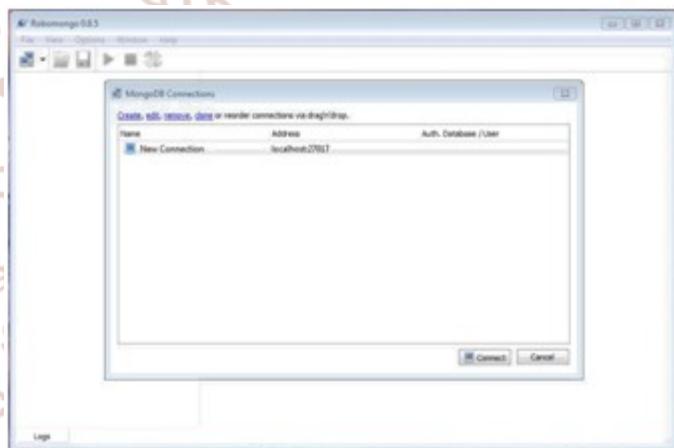


Figure4: MongoDB connection to Tomcat server



Figure5: Soldier Login page

The Soldier has to register if he is a new user. The figure 5 shows the login UI for soldier. The soldier has to enter the valid registered credentials like Email-id and password to upload the text and unstructured data to cloud. On click of submit button soldier will be navigated to user menus page.

**Figure 6: Actions performed by Soldier**

The figure 6 shows the list of actions performed by soldier such as upload-daily-data, upload plan and view reply. On click of upload-daily-data link the soldier will be navigated to upload details page and is able to upload data.

figure7: Soldier uploads data by entering required, inputs

The figure 7 shows the screen shot of uploading data by soldier. Soldier has to enter secret-id, current location, situation of current location and he has to report. On click of submit button data will be successfully uploaded.

Figure 8: Commander Login page

The commander has to register if he is new user. The figure 8 shows the login UI for commander. The commander has to enter the valid registered credentials like email-id and password to view the text and unstructured data. On click of submit button commander will be navigated to commander menus page.

ID	Email	Value	Time	Status
101	user1@gmail.com	101	2013-08-01 10:00:00	status1
102	user2@gmail.com	102	2013-08-01 10:00:00	status2
103	user3@gmail.com	103	2013-08-01 10:00:00	status3
104	user4@gmail.com	104	2013-08-01 10:00:00	status4
105	user5@gmail.com	105	2013-08-01 10:00:00	status5
106	user6@gmail.com	106	2013-08-01 10:00:00	status6
107	user7@gmail.com	107	2013-08-01 10:00:00	status7
108	user8@gmail.com	108	2013-08-01 10:00:00	status8
109	user9@gmail.com	109	2013-08-01 10:00:00	status9
110	user10@gmail.com	110	2013-08-01 10:00:00	status10
111	user11@gmail.com	111	2013-08-01 10:00:00	status11
112	user12@gmail.com	112	2013-08-01 10:00:00	status12
113	user13@gmail.com	113	2013-08-01 10:00:00	status13
114	user14@gmail.com	114	2013-08-01 10:00:00	status14
115	user15@gmail.com	115	2013-08-01 10:00:00	status15
116	user16@gmail.com	116	2013-08-01 10:00:00	status16
117	user17@gmail.com	117	2013-08-01 10:00:00	status17
118	user18@gmail.com	118	2013-08-01 10:00:00	status18
119	user19@gmail.com	119	2013-08-01 10:00:00	status19
120	user20@gmail.com	120	2013-08-01 10:00:00	status20
121	user21@gmail.com	121	2013-08-01 10:00:00	status21
122	user22@gmail.com	122	2013-08-01 10:00:00	status22
123	user23@gmail.com	123	2013-08-01 10:00:00	status23
124	user24@gmail.com	124	2013-08-01 10:00:00	status24
125	user25@gmail.com	125	2013-08-01 10:00:00	status25
126	user26@gmail.com	126	2013-08-01 10:00:00	status26
127	user27@gmail.com	127	2013-08-01 10:00:00	status27
128	user28@gmail.com	128	2013-08-01 10:00:00	status28
129	user29@gmail.com	129	2013-08-01 10:00:00	status29
130	user30@gmail.com	130	2013-08-01 10:00:00	status30
131	user31@gmail.com	131	2013-08-01 10:00:00	status31
132	user32@gmail.com	132	2013-08-01 10:00:00	status32
133	user33@gmail.com	133	2013-08-01 10:00:00	status33
134	user34@gmail.com	134	2013-08-01 10:00:00	status34
135	user35@gmail.com	135	2013-08-01 10:00:00	status35
136	user36@gmail.com	136	2013-08-01 10:00:00	status36
137	user37@gmail.com	137	2013-08-01 10:00:00	status37
138	user38@gmail.com	138	2013-08-01 10:00:00	status38
139	user39@gmail.com	139	2013-08-01 10:00:00	status39
140	user40@gmail.com	140	2013-08-01 10:00:00	status40
141	user41@gmail.com	141	2013-08-01 10:00:00	status41
142	user43@gmail.com	142	2013-08-01 10:00:00	status42
143	user44@gmail.com	143	2013-08-01 10:00:00	status43
144	user45@gmail.com	144	2013-08-01 10:00:00	status44
145	user46@gmail.com	145	2013-08-01 10:00:00	status45
146	user47@gmail.com	146	2013-08-01 10:00:00	status46
147	user48@gmail.com	147	2013-08-01 10:00:00	status47
148	user49@gmail.com	148	2013-08-01 10:00:00	status48
149	user50@gmail.com	149	2013-08-01 10:00:00	status49
150	user51@gmail.com	150	2013-08-01 10:00:00	status50
151	user52@gmail.com	151	2013-08-01 10:00:00	status51
152	user53@gmail.com	152	2013-08-01 10:00:00	status52
153	user54@gmail.com	153	2013-08-01 10:00:00	status53
154	user55@gmail.com	154	2013-08-01 10:00:00	status54
155	user56@gmail.com	155	2013-08-01 10:00:00	status55
156	user57@gmail.com	156	2013-08-01 10:00:00	status56
157	user58@gmail.com	157	2013-08-01 10:00:00	status57
158	user59@gmail.com	158	2013-08-01 10:00:00	status58
159	user60@gmail.com	159	2013-08-01 10:00:00	status59
160	user61@gmail.com	160	2013-08-01 10:00:00	status60
161	user62@gmail.com	161	2013-08-01 10:00:00	status61
162	user63@gmail.com	162	2013-08-01 10:00:00	status62
163	user64@gmail.com	163	2013-08-01 10:00:00	status63
164	user65@gmail.com	164	2013-08-01 10:00:00	status64
165	user66@gmail.com	165	2013-08-01 10:00:00	status65
166	user67@gmail.com	166	2013-08-01 10:00:00	status66
167	user68@gmail.com	167	2013-08-01 10:00:00	status67
168	user69@gmail.com	168	2013-08-01 10:00:00	status68
169	user70@gmail.com	169	2013-08-01 10:00:00	status69
170	user71@gmail.com	170	2013-08-01 10:00:00	status70
171	user72@gmail.com	171	2013-08-01 10:00:00	status71
172	user73@gmail.com	172	2013-08-01 10:00:00	status72
173	user74@gmail.com	173	2013-08-01 10:00:00	status73
174	user75@gmail.com	174	2013-08-01 10:00:00	status74
175	user76@gmail.com	175	2013-08-01 10:00:00	status75
176	user77@gmail.com	176	2013-08-01 10:00:00	status76
177	user78@gmail.com	177	2013-08-01 10:00:00	status77
178	user79@gmail.com	178	2013-08-01 10:00:00	status78
179	user80@gmail.com	179	2013-08-01 10:00:00	status79
180	user81@gmail.com	180	2013-08-01 10:00:00	status80
181	user82@gmail.com	181	2013-08-01 10:00:00	status81
182	user83@gmail.com	182	2013-08-01 10:00:00	status82
183	user84@gmail.com	183	2013-08-01 10:00:00	status83
184	user85@gmail.com	184	2013-08-01 10:00:00	status84
185	user86@gmail.com	185	2013-08-01 10:00:00	status85
186	user87@gmail.com	186	2013-08-01 10:00:00	status86
187	user88@gmail.com	187	2013-08-01 10:00:00	status87
188	user89@gmail.com	188	2013-08-01 10:00:00	status88
189	user90@gmail.com	189	2013-08-01 10:00:00	status89
190	user91@gmail.com	190	2013-08-01 10:00:00	status90
191	user92@gmail.com	191	2013-08-01 10:00:00	status91
192	user93@gmail.com	192	2013-08-01 10:00:00	status92
193	user94@gmail.com	193	2013-08-01 10:00:00	status93
194	user95@gmail.com	194	2013-08-01 10:00:00	status94
195	user96@gmail.com	195	2013-08-01 10:00:00	status95
196	user97@gmail.com	196	2013-08-01 10:00:00	status96
197	user98@gmail.com	197	2013-08-01 10:00:00	status97
198	user99@gmail.com	198	2013-08-01 10:00:00	status98
199	user100@gmail.com	199	2013-08-01 10:00:00	status99
200	user101@gmail.com	200	2013-08-01 10:00:00	status100
201	user102@gmail.com	201	2013-08-01 10:00:00	status101
202	user103@gmail.com	202	2013-08-01 10:00:00	status102
203	user104@gmail.com	203	2013-08-01 10:00:00	status103
204	user105@gmail.com	204	2013-08-01 10:00:00	status104
205	user106@gmail.com	205	2013-08-01 10:00:00	status105
206	user107@gmail.com	206	2013-08-01 10:00:00	status106
207	user108@gmail.com	207	2013-08-01 10:00:00	status107
208	user109@gmail.com	208	2013-08-01 10:00:00	status108
209	user110@gmail.com	209	2013-08-01 10:00:00	status109
210	user111@gmail.com	210	2013-08-01 10:00:00	status110
211	user112@gmail.com	211	2013-08-01 10:00:00	status111
212	user113@gmail.com	212	2013-08-01 10:00:00	status112
213	user114@gmail.com	213	2013-08-01 10:00:00	status113
214	user115@gmail.com	214	2013-08-01 10:00:00	status114
215	user116@gmail.com	215	2013-08-01 10:00:00	status115
216	user117@gmail.com	216	2013-08-01 10:00:00	status116
217	user118@gmail.com	217	2013-08-01 10:00:00	status117
218	user119@gmail.com	218	2013-08-01 10:00:00	status118
219	user120@gmail.com	219	2013-08-01 10:00:00	status119
220	user121@gmail.com	220	2013-08-01 10:00:00	status120
221	user122@gmail.com	221	2013-08-01 10:00:00	status121
222	user123@gmail.com	222	2013-08-01 10:00:00	status122
223	user124@gmail.com	223	2013-08-01 10:00:00	status123
224	user125@gmail.com	224	2013-08-01 10:00:00	status124
225	user126@gmail.com	225	2013-08-01 10:00:00	status125
226	user127@gmail.com	226	2013-08-01 10:00:00	status126
227	user128@gmail.com	227	2013-08-01 10:00:00	status127
228	user129@gmail.com	228	2013-08-01 10:00:00	status128
229	user130@gmail.com	229	2013-08-01 10:00:00	status129
230	user131@gmail.com	230	2013-08-01 10:00:00	status130
231	user132@gmail.com	231	2013-08-01 10:00:00	status131
232	user133@gmail.com	232	2013-08-01 10:00:00	status132
233	user134@gmail.com	233	2013-08-01 10:00:00	status133
234	user135@gmail.com	234	2013-08-01 10:00:00	status134
235	user136@gmail.com	235	2013-08-01 10:00:00	status135
236	user137@gmail.com	236	2013-08-01 10:00:00	status136
237	user138@gmail.com	237	2013-08-01 10:00:00	status137
238	user139@gmail.com	238	2013-08-01 10:00:00	status138
239	user140@gmail.com	239	2013-08-01 10:00:00	status139
240	user141@gmail.com	240	2013-08-01 10:00:00	status140
241	user142@gmail.com	241	2013-08-01 10:00:00	status141
242	user143@gmail.com	242	2013-08-01 10:00:00	status142
243	user144@gmail.com	243	2013-08-01 10:00:00	status143
244	user145@gmail.com	244	2013-08-01 10:00:00	status144
245	user146@gmail.com	245	2013-08-01 10:00:00	status145
246	user147@gmail.com	246	2013-08-01 10:00:00	status146
247	user148@gmail.com	247	2013-08-01 10:00:00	status147
248	user149@gmail.com	248	2013-08-01 10:00:00	status148
249	user150@gmail.com	249	2013-08-01 10:00:00	status149
250	user151@gmail.com	250	2013-08-01 10:00:00	status150
251	user152@gmail.com	251	2013-08-01 10:00:00	status151
252	user153@gmail.com	252	2013-08-01 10:00:00	status152
253	user154@gmail.com	253	2013-08-01 10:00:00	status153
254	user155@gmail.com	254	2013-08-01 10:00:00	status154
255	user156@gmail.com	255	2013-08-01 10:00:00	status155
256	user157@gmail.com	256	2013-08-01 10:00:00	status156
257	user158@gmail.com	257	2013-08-01 10:00:00	status157
258	user159@gmail.com	258	2013-08-01 10:00:00	status158
259	user160@gmail.com	259	2013-08-01 10:00:00	status159
260	user161@gmail.com	260	2013-08-01 10:00:00	status160
261	user162@gmail.com	261	2013-08-01 10:00:00	status161
262	user163@gmail.com	262	2013-08-01 10:00:00	status162
263	user164@gmail.com	263	2013-08-01 10:00:00	status163
264	user165@gmail.com	264	2013-08-01 10:00:00	status164
265	user166@gmail.com	26		

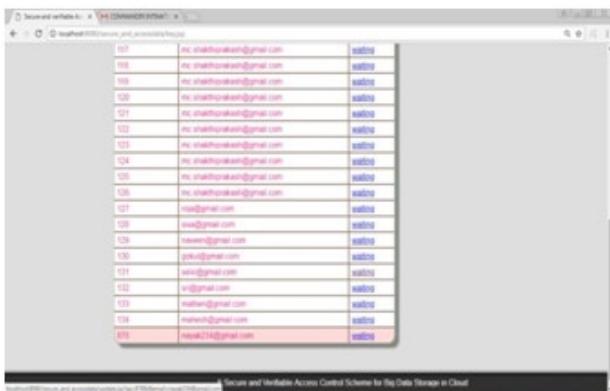


Figure 12: Screenshot of view-daily-data table in Cloud

The figure 12 shows the screenshot of view-daily-data screen. On click of waiting link the generated key will be sent to corresponding commander's email-id.

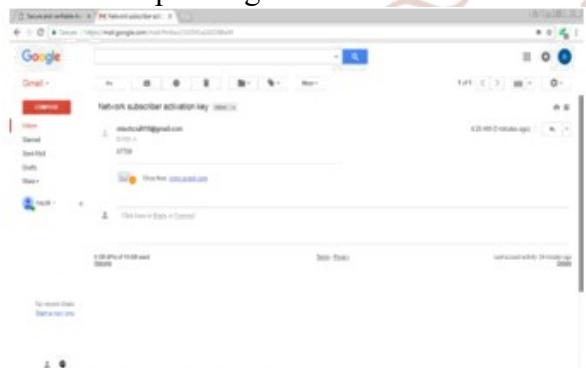


Figure 13: Screenshot of Commander’s mail-id

Figure 14: Screenshot of decrypted data

Figure 10. A screenshot of decrypted data						
	username	password	last login	status	roles	Ready
manish	manish@gmail.com	154	20/09/17 12:00am	status	Videos	Ready
wilma	wilma@gmail.com	155	20/09/17 12:00am	status	Videos	Ready
vivek	vivek@net.com	156	20/09/17 12:00am	status	Videos	Ready
vinay	vinay@gmail.com	157	20/09/17 12:00am	status	Videos	Ready
venu	venu@gmail.com	158	20/09/17 12:00am	status	Videos	Ready
viraj	viraj@gmail.com	159	20/09/17 12:00am	status	Videos	Ready
virajendra	virajendra123@gmail.com	160	20/09/17 12:00am	status	Videos	Ready
devash	devashreyekar@gmail.com	161	20/09/17 12:00am	status	Videos	Ready
suparna	suparna123@gmail.com	162	4/04/17 12:00pm	status	Videos	Ready
praveen	praveen12345@gmail.com	163	04/04/17 12:00pm	status	Videos	Ready
seema	seema@varun123@gmail.com	164	21/03/1998 11:59pm	status	Videos	Ready
flip	seema@varun123@gmail.com	1654	21/03/1998 11:59pm	status	Videos	Ready
manu	seema@varun123@gmail.com	1658	23/04/1998 11:59pm	status	Videos	Ready
flip	seema@varun123@gmail.com	1654	9/10/2016 11:59pm	status	Videos	Ready
flip	seema@varun123@gmail.com	1657	22/07/2016 11:59pm	status	Videos	Ready
abid	abid@jpmail.com	167	17/03/2010 11:59pm	status	Videos	Ready
Naresh	naresh123@gmail.com	179	23/07/2009 12:59pm	status	Videos	Ready
Naresh	naresh123@gmail.com	187	23/07/2009 12:59pm	status	Videos	Ready

Figure 15: Screenshot of Reply option for commander

The figure 15 shows the screenshot of reply option for commander. On click of reply link, the commander can send reply to soldier.



Figure 16: Screenshot of view reply page

VIII. CONCLUSION AND FUTURE SCOPE

The projected framework was created along with an intention of giving thorough protection toward the secret data when distributed via the cloud, it regularly inclined to protection dangers by means of the interlopers with the hidden NTRU positioned encryption and decryption strategies. Evidently, it discovers a tremendous request in the section of Intelligence and digital dangers in different inward and outer resistance framework. Henceforth, the projected framework reproduces the run of the mill data distributing procedure amid the resistance work force and effectively demonstrates that the information distributed by means of the cloud is enveloped along with fundamental protection highlights and available just the approved event. The projected framework has thoroughly dissected the rightness, protection quality, and calculation multifaceted nature of the projected conspire.

In spite of the fact that the framework gives sufficient protection, however the likelihood of furious the protection stays less, the ascent in intense calculation equipment that utilization dimension processing procedures may at present have the capacity to disentangle the data effectively. Thus, a calculation that can identify the gatecrasher's action in view of authentic records would be alluring in outlook and henceforth an artificial engineering learning calculations for abnormality identification may be conveyed over the current arrangement. The outlook investigation incorporates enhancing the plan through joining limit mystery imparting to trait positioned admission manage, it includes an entrance arrangement that could put different prerequisites for

a client to unscramble a sub contracted figure content statistics located at the cloud.

The future research includes improving the scheme by combining threshold secret sharing with attribute based access control, which involves an access structure that can place various requirements for a user to decrypt an outsourced cipher text data in the cloud.

IX. REFERENCES

1. Chunqiang Hu et al., "A Secure and Verifiable Access Control Scheme for Big Data Storage in Clouds" IEEE Transactions on Big data, Year: 2017, Volume: PP, Issue: 99
2. C. Hu, H. Li, Y. Huo, T. Xiang, and X. Liao, "Secure and efficient data communication protocol for wireless body area networks," IEEE Transactions on Multi-Scale Computing Systems, vol. 2, no. 2, pp. 94–107, 2016.
3. C. Hu, X. Cheng, Z. Tian, J. Yu, K. Akkaya, and L. Sun, "An attribute based signcryption scheme to secure attribute-defined multicast communications," in SecureComm 2015. Springer, 2015, pp. 418–435.
4. C. Hu, N. Zhang, H. Li, X. Cheng, and X. Liao, "Body area network security: a fuzzy attribute-based signcryption scheme," IEEE journal on selected areas in communications, vol. 31, no. 9, pp. 37–46, 2013.
5. V. Marx, "Biology: The big challenges of big data," Nature, vol. 498, no. 7453, pp. 255–260, 2013.
6. C. Hu, F. Zhang, X. Cheng, X. Liao, and D. Chen, "Securing communications between external users and wireless body area networks," in Proceedings of the 2nd ACM workshop on Hot topics on wireless network security and privacy. ACM, 2013, pp. 31–36.
7. M. A. Beyer and D. Laney, "The importance of big data: a definition," Stamford, CT: Gartner, 2012.
8. B. Waters, "Cipher text-policy attribute-based encryption: An expressive, efficient, and provably secure realization," Public Key Cryptography–PKC 2011, pp. 53–70, 2011.
9. J. Hoffstein, "NTRU Public Key Cryptosystem–Methodology", shodhganga.inflibnet.ac.in/ bit stream/1060/3/103254/10/10_chapter-iii.pdf.